

Divide and Conquer to Multiply and Order

Reading: Chapter 18

Divide-and-conquer is a frequently-useful algorithmic technique tied up in recursion.

We'll see how it is useful in

- SORTING
- MULTIPLICATION

A divide-and-conquer algorithm has three basic steps...

1. *Divide problem into smaller versions of the same problem.*
2. *Recursively solve each smaller version.*
3. *Combine solutions to get overall solution.*

Sorting

Problem SORT:

Input: *an array A of integers.*

Output: *an array with all elements of A in increasing order.*

example:

```

+-----+-----+-----+-----+-----+-----+-----+-----+
A = | 19 | 1 | 29 | 30 | 6 | 15 | 2 | 5 |
+-----+-----+-----+-----+-----+-----+-----+-----+

+-----+-----+-----+-----+-----+-----+-----+-----+
return | 1 | 2 | 5 | 6 | 15 | 19 | 29 | 30 |
+-----+-----+-----+-----+-----+-----+-----+-----+

```

Merge-Sort

Algorithm $M\text{-Sort}(A)$:

if A 's length > 1 , then:

return A .

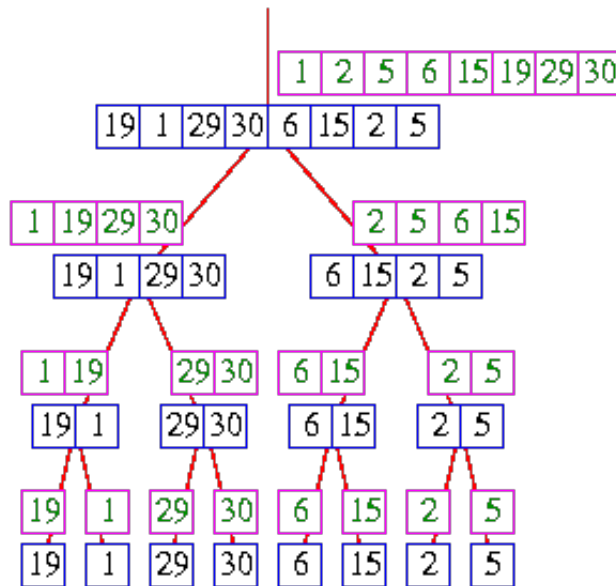
else:

Let A_L be the first half of A .

Let A_R be the second half of A .

return $Merge(M\text{-Sort}(A_L), M\text{-Sort}(A_R))$

end of if



Merging

Algorithm Merge(A, B):

Let n_A and n_B be the length of arrays A and B.

// C holds merged array

Let C be an array of length $n_A + n_B$.

// a,b,c are current positions in A,B,C

Let a, b, and c hold 0.

while $a < n_A$ **and** $b < n_B$, **do:**

if $a < n_A$ **and** $A[a] < B[b]$, **then:**

 Let C[c] hold A[a].

 Add 1 to a and c.

else:

 Let C[c] hold B[b].

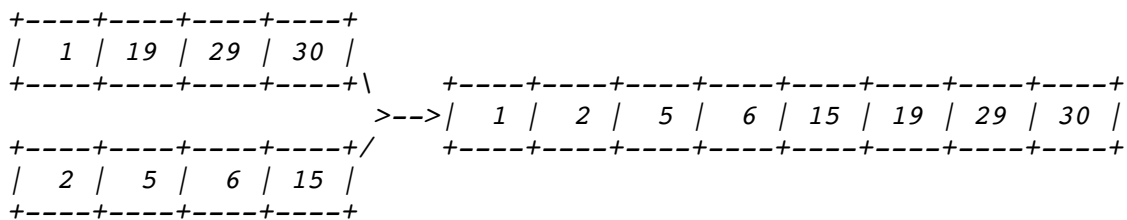
 Add 1 to b and c.

end of if

end of loop

return C.

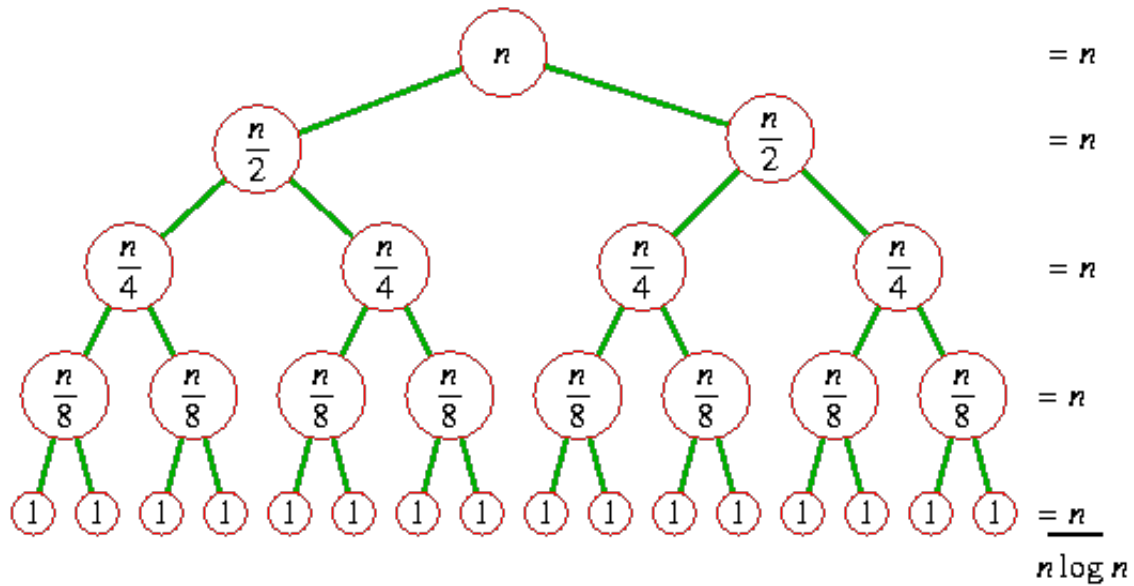
example:



How Much Time?

Write a recurrence and solve...

$$T(n) = 2 T(n / 2) + cn$$



Multiplication

Problem MULTIPLICATION

Input: two n -digit integers a and b

Output: product of a and b

example:

```

      1980 = a
x     2315 = b
-----
      9900
      1980
      5940
+   3960
-----
    4573700 = a x b
    
```

This is the algorithm you learned in grade school. Notice it takes $O(n^2)$ time.

Dividing the Problem

We divide each integer into two halves.

$$\begin{array}{l}
 aL = 19 \quad | \quad 80 = aR \\
 bL = 23 \quad | \quad 15 = bR
 \end{array}$$

$$\begin{array}{r}
 \quad \quad \quad aL \quad \quad \quad aR \\
 x \quad \quad \quad bL \quad \quad \quad bR \\
 \hline
 \quad \quad \quad aL \ bR \quad \quad \quad aR \ bR
 \end{array}$$

$$\begin{array}{r}
 + \quad a_L \quad b_L \qquad \qquad a_R \quad b_L \\
 \hline
 a_L \quad b_L \quad a_L \quad b_R + a_R \quad b_L \quad a_R \quad b_R
 \end{array}$$

So our algorithm is to compute $a_L b_L$, $a_L b_R$, $a_R b_L$, and $a_R b_R$, and add.

$$T(n) \leq 4 T(n / 2) + O(n)$$

$$T(n) = O(n^2)$$

Algorithm Divide-Mult(a,b):

if a or b has one digit, **then:**

 return $a * b$.

else:

 Let n be the number of digits in $\max\{a, b\}$.

 Let a_L and a_R be left and right halves of a .

 Let b_L and b_R be left and right halves of b .

 Let x_1 hold Divide-Mult(a_L, b_L).

 Let x_2 hold Divide-Mult(a_L, b_R).

 Let x_3 hold Divide-Mult(a_R, b_L).

 Let x_4 hold Divide-Mult(a_R, b_R).

return $x_1 * 10^n + (x_2 + x_3) * 10^{n/2} + x_4$.

end of if

Being Clever

We can actually get away with just three multiplications!

$$\begin{aligned}
 x_1 &= a_L b_L \\
 x_2 &= a_R b_R \\
 x_3 &= (a_L + a_R) (b_L + b_R)
 \end{aligned}$$

$$\begin{array}{r}
 \qquad \qquad \qquad a_L \qquad \qquad a_R \\
 x \qquad \qquad \qquad b_L \qquad \qquad b_R \\
 \hline
 a_L b_L \quad a_L b_R + a_R b_L \quad a_R b_R \\
 x_1 \qquad \quad x_3 - x_1 - x_2 \qquad x_2
 \end{array}$$

Now our algorithm is to compute x_1 , x_2 , and x_3 , find $x_3 - x_1 - x_2$, and add.

Algorithm Karatsuba(a,b):

if a or b has one digit, **then:**

 return $a * b$.

else:

 Let n be the number of digits in $\max\{a, b\}$.

 Let a_L and a_R be left and right halves of a .

 Let b_L and b_R be left and right halves of b .

 Let x_1 hold Karatsuba(a_L, b_L).

 Let x_2 hold Karatsuba($a_L + a_R, b_L + b_R$).

 Let x_3 hold Karatsuba(a_R, b_R).

return $x_1 * 10^n + (x_2 - x_1 - x_3) * 10^{n/2} + x_3$.

end of if

 Ooooh!

 Aaaah!

An Example

```

IN Multiply(1980, 2315)
  19 80
  23 15
    IN Multiply(19, 23)
      1 9
      2 3
        IN Multiply(1, 2)
          return 2
        IN Multiply(9, 3)
          return 27
        IN Multiply(10, 5)
          1 0
          0 5
            IN Multiply(1, 0)
              return 0
            IN Multiply(0, 5)
              return 0
            IN Multiply(1, 5)
              return 5
          5 - 0 - 0 = 5
          0
          5
          0
          ---
          50
        return 50
      50 - 27 - 2 = 21
      2
      21
      27
      ---
      437
    return 437
  IN Multiply(80, 15)
    8 0
    1 5
      IN Multiply(8, 1)
        return 8
      IN Multiply(0, 5)
        return 0
      IN Multiply(8, 6)
        return 48
    48 - 8 - 0 = 40
    8
    40
    0
    ----
    1200
  return 1200

```

```

IN Multiply(99, 38)
 9 9
 3 8
  IN Multiply(9, 3)
  return 27
  IN Multiply(9, 8)
  return 72
  IN Multiply(18, 11)
  1 8
  1 1
    IN Multiply(1, 1)
    return 1
    IN Multiply(8, 1)
    return 8
    IN Multiply(9, 2)
    return 18
  18 - 8 - 1 = 9
  1
  9
  8
  ---
  198
  return 198
198 - 27 - 72 = 99
27
 99
 72
----
3762
return 3762
3762 - 437 - 1200 = 2125

437
2125
 1200
-----
4583700
return 4583700

```

How Do Multiplication Algorithms Compare in Practice?

